

Recursion

Data structures and Algorithms



What is recursion?

- ◆ A way of thinking about problems.
- ◆ A method for solving problems.
- ◆ Related to mathematical induction.

- ◆ In programming:

- a function calls itself
 - ◆ direct recursion

```
int f () {  
    ... f(); ...  
}
```

- a function calls its invoker
 - ◆ indirect recursion

```
int f () {  
    ... g(); ...  
}
```

```
int g () {  
    ... f(); ...  
}
```

Outline of a Recursive Function

FIGURE 7.1

A Set of Nested Wooden Figures



base case →

if (answer is known)
 provide the answer
else

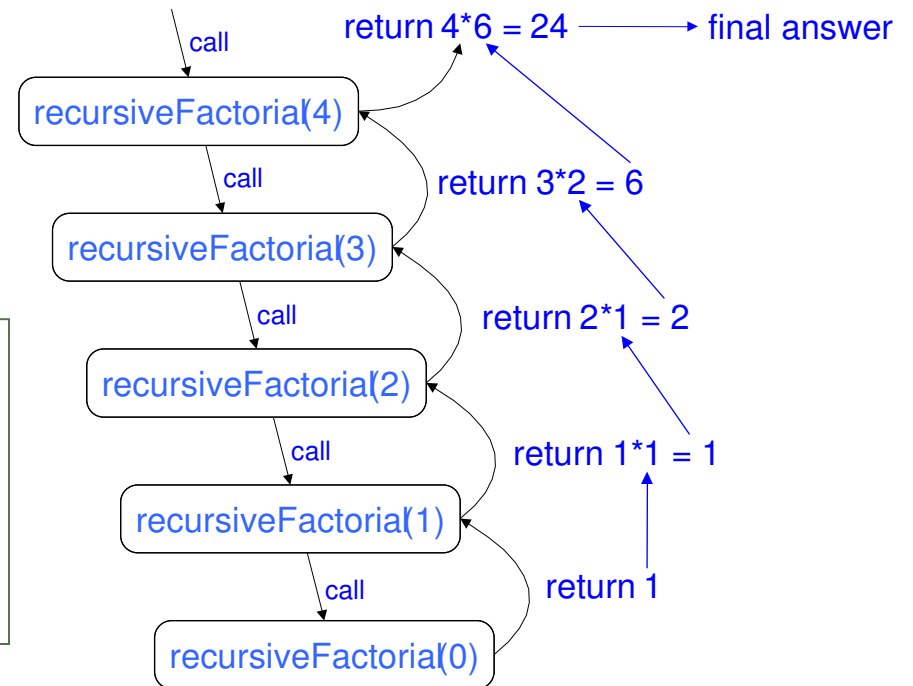
recursive case →

 make a recursive call
 to solve a **smaller** version
 of the **same** problem

Recursive Factorial Method

- ◆ $n! = n * (n-1) * (n-2) * \dots * 3 * 2 * 1$
- ◆ $n! = n * (n-1)!$
- ◆ $0! = 1$

```
Algorithm recursiveFactorial(n)  
  if  $n == 0$  then  
    return 1  
  else  
    return  $n * \text{recursiveFactorial}(n-1)$ 
```



Fibonacci sequence

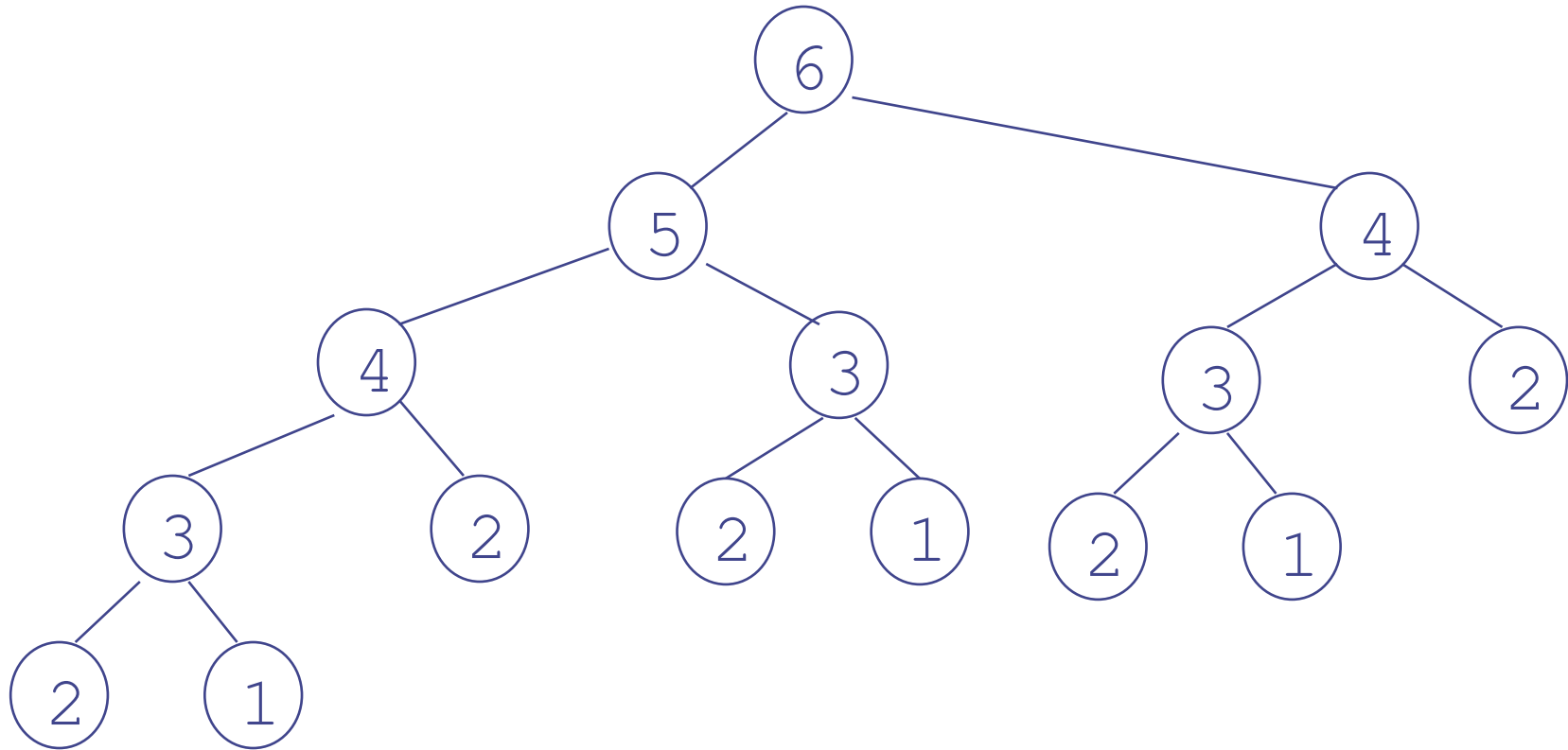
1, 1, 2, 3, 5, 8, 13, 21,

$$\text{fib}(n) = \begin{cases} 1 & \text{for } n == 1 \\ 1 & \text{for } n == 2 \\ \text{fib}(n-2) + \text{fib}(n-1) & \text{for } n > 2 \end{cases}$$

```
Algorithm fib(n)  
  if n ≤ 2 then  
    return 1  
  else  
    return fib(n-2) + fib(n-1)
```

Tracing fib(6)

```
Algorithm fib(n)  
  if n ≤ 2 then  
    return 1  
  else  
    return fib(n-2) + fib(n-1)
```



computation repeats!

Design a Recursive Algorithm

- ◆ There must be at least one case (the base case), for a small value of n , that can be solved directly
- ◆ A problem of a given size n can be split into one or more *smaller* versions of the *same* problem (recursive case)
- ◆ Recognize the base case and provide a solution to it
- ◆ Devise a strategy to split the problem into smaller versions of itself while making progress toward the base case
- ◆ Combine the solutions of the smaller problems in such a way as to solve the larger problem

Euclid's Algorithm

- ◆ Finds the greatest common divisor of two nonnegative integers that are not both 0
- ◆ Recursive definition of gcd algorithm
 - $\text{gcd}(a, b) = a$ (if b is 0)
 - $\text{gcd}(a, b) = \text{gcd}(b, a \% b)$ (if $b \neq 0$)
- ◆ Implementation:

```
int gcd (int a, int b)
{
    if (b == 0)
        return a;
    else
        return gcd (b, a % b);
}
```


Iterative vs. recursive gcd

```
int gcd (int a, int b)
{
    int temp;
    while (b != 0)
    {
        temp = b;
        b = a % b;
        a = temp;
    }
    return a;
}
```

```
int gcd (int a, int b)
{
    if (b == 0)
        return a;
    else
        return gcd (b, a % b);
}
```

Multiple recursion

- ◆ Tail recursion: a linearly recursive method makes its recursive call as its last step.
 - e.g. recursive gcd
 - Can be easily converted to non-recursive methods
- ◆ Binary recursion: there are **two** recursive calls for each non-base case
 - e.g. fibonacci sequence
- ◆ Multiple recursion: makes potentially **many** recursive calls (not just one or two).

Multiple recursion – Example

List all 'abc' strings of length l

```
void listAllStrings(int length, char* start)
{
    if (length < 1) { //base case: empty string
        *start = '\0';
        output();
    } else { //recursive case: reduce length by 1
        for (char c = 'a'; c <= 'c'; c++) {
            *start = c;
            listAllStrings(length-1, start+1);
        }
    }
}
```

aaa
aab
aac
aba
abb
abc
aca
acb
acc
baa
bab
...
ccc

Why using recursion?

◆ Recursion makes your code **faster? No!**

- overhead for function call and return
- values recomputed

◆ Recursion uses **less memory? No!**

- overhead for a function call and return (stack memory)

◆ Recursion makes your code **simple? Sometimes.**

- readable code that is easy to debug